PetriScript Reference Manual 1.0

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Introduction

The CPN-AMI platform provides many tools to work on Petri nets such as verifying or model-checking tools. It was easily possible to graphically design simple Petri nets with Macao¹, but various works made internally at LIP6 reveal that it was needed to automate such task.

Therefore PetriScript has been designed to provide some facilities in modelling placestransition and coloured Petri nets within the CPN-AMI platform².

Its main purpose is to automate modelling operations on Petri nets such as merging, creating, and connecting nodes. Thus, it supports almost everything needed like macros, loops control, lists, string and arithmetic expressions, ..., and avoids to the maximum the intervention of the user. Its syntax is more or less Ada-like.

Chapter 1 will introduce the basic aspects of PetriScript such as loop, variables, ..., while chapter 2 will explain the manipulation of Petri nets via PetriScript .

A Backus-Naur Form is provided in appendix A.

¹www-src.lip6.fr/logiciels/mars/MACAO

²www-src.lip6.fr/logiciels/mars/CPNAMI

Chapter 1

PetriScript basics

Before learning the manipulation of Petri nets with PetriScript , which is presented in chapter 2, there are some aspects of the language to understand.

1.1 Using PetriScript

To launch PetriScript in Macao, select AMI-NET > Modelling facilities > Petri Net assembling > Execute list of commands. Then Macao opens a window, shown in figure 1.1.



Figure 1.1: PetriScript window

It is strongly advised to type the program in a real text editor like Emacs¹, or at least

¹PetriScript provides you a mode for Emacs

to copy your script before validating in order to keep the text in case of an error.

You can also choose to launch PetriScript scripts in a debug mode by selecting AMI-NET > Modelling Facilites > Petri Net assembling > Execute list of commands (debug mode). This time, if an error occurs, then Macao will display state in which state the net just was before the error happened.

1.2 Generalities

Typically, PetriScript consists in sequence of declarations followed by a sequence of instructions. It's also possible to define some macros anywhere in the text. Declarations, instructions and macros definitions all end with the symbol ';'. There can be as many spaces and carriage return as you wish. Like Ada, it is verbose, but unlike Ada, it is case-sensitive to preserve names and attributes of nodes.

1.3 Comments

Comments begin with '--' and finish with a carriage return.

1.4 The basic instruction: print

It prints the following string expression (see 1.8) into the historic window of Macao. So, a first very simple script would be like this:

```
print "Hello world!";
```

1.5 Macros

PetriScript uses m4² as its preprocessor. So a macro is defined this way:

```
define(macro, text)
```

It has the effect to replace all 'macro' which follow this definition in the script by 'text'. So the tiny following scripts produce "Hello World!" on the historic window of Macao:

```
define(HELLO, Hello World!)
print "HELLO";
define(HELLO, "Hello World!")
print HELLO;
```

You can use any facilities of m4 such as testing if a macro is defined. For more details, take a look at the m4 documentation.

1.6 Variables types

Three different types are available in PetriScript: integers, lists and strings. Names of variables which hold these types always begin with the symbol '\$' and can contain numerical characters, letters and the underscore symbol '_'. It is completely legal to write \$a_very_long_variable_name_993023.

http://www.gnu.org/software/m4/m4.html

Variables can be affected with a value or with another variable of the same type at declaration or anywhere in the script using ':='.

Declarations are made at the beginning of the script before any instruction. Type is given before the variable name and an affectation if needed.

Example:

```
int $i;
string $chaine := "a little text";
```

Integers

They are declared by the keyword int.

```
\begin{array}{ll} \mbox{int $\$i$;} \\ \mbox{int $\$j$} := 10; \\ \end{array}
```

Any arithmetic expression (see 1.7) can be affected to an integer variable.

```
i := 1 + 10 \mod 2;
i := i / 3;
```

Lists

They are declared by the keyword list.

```
list $1;
list $m := place {"a" , "b"};
```

Elements of the list are separated by commas and surrounded by brackets. As their behaviour is completely related to the manipulation of nodes, the exact way of using these lists is described further on, at section 2.1.2, p. 11.

• Strings

They are declared by the keyword string.

```
string $s;
string $t := "some text to print";
```

Any string expression (see 1.8) can be affected to a string variable.

```
$s := "text" & '10/2';
$t := '$s' & "an other text";
```

1.7 Arithmetic expressions

PetriScript handles arithmetical operations on integers. The available operations are:

- multiplication '*'
- division '/'
- modulo 'mod'
- subtraction '-'
- addition '+'

Operands of an arithmetic expression can be integers or variables.

As usual, multiplication, division and modulo operators have the same priority. They have priority over addition and subtraction. Obviously, if you have to deal with complex arithmetic expressions, you can use parentheses.

```
i := 1+2*3;

i := i \mod 3 * (2 / 5);
```

If a list is given in an arithmetic expression, it returns its size. It is not legal to give a string expression.

1.8 String expressions

String expressions are used by the **print** instruction and to describe nodes (explained in chapter 2).

A string expression consists in the concatenation of different operands with the operator '&'. An operand can be a

• sequence of characters surrounded by "":

```
"Some text"
```

• variable surrounded by ''':

```
'$variable'
```

If the given variable is a string, it just returns its content; when it is an integer, it returns its value interpolated into a string. Finally, if the variable is a list, it returns a string representation of all its components.

arithmetic expression surrounded by ''':

```
'10 / 5 +3'
'2 + $i'
```

So, the following script will produce "Result: 5 + 2" on the historic window.

```
int $i;
string $s;

$i := 10 / 2;
$s := "Result: " & '$i' & " + " & '10 mod 8';

print '$s';
```

To insert a carriage return in a string expression, use the $^\prime/^\prime$ character: "a line $^\prime$ another line"

1.9 Control structures

PetriScript provides the classic ways of controlling your scripts with conditional and loop structures.

1.9.1 Conditional structures

```
if boolean_expression then
    instructions_sequence
end if

if boolean_expression then
    instructions_sequence
else
    instructions_sequence
end if
```

A *boolean_expression* is composed of arithmetic expressions connected with boolean operators, showed in table 1.1.

Symbol	Comparison		
=	Equal		
/=	Different		
<=	Inferior or equal		
>=	Superior or equal		
<	Inferior		
>	Superior		

Table 1.1: Boolean Operators

A boolean expression can be composed of two boolean expressions connected with **and** and **or**. It can also be negated with **not**.

Examples:

You can also test if a value is in a given range with the boolean expression *arithmetic* expression *in* arithmetic expression .. arithmetic expression.

Examples:

1.9.2 Loop structures

You can use two types of structures to loop, one which loops while the given boolean expression is verified, and another one which loops for a given range.

• While structure

```
while boolean_expression loop
instructions_sequence
end loop
```

Example:

• For structure

```
for variable in arithmetic_expression..arithmetic_expression loop
    instructions_sequence
end loop
```

Example:

1.10 Basic example

Here is a little script which presents every basic aspect of PetriScript:

CHAPTER 1. PETRISCRIPT BASICS

```
$s := '$s' & '$i' & EVEN;
else
$s := '$s' & '$i' & ODD;
end if;
end loop;
print "Result: " & '$s';
```

It displays "Result: 1 odd. 2 even. 3 odd. 4 even. 5 odd. 6 even. 7 odd. 8 even. 9 odd. 10 even." on the historic window.

Chapter 2

Manipulate Petri Nets

This chapter explains how to manipulate Petri nets with PetriScript.

2.1 Describe a node

Before acting on Petri nets, it is necessary to describe its components: nodes and arcs.

2.1.1 Single node

A node is designated by its name, which is a string expression and its type. The type of a node can be a **place**, a **queue**, an **immediate** or a **transition**. The form to designate a node is the following: *type node_name*. For example, the place *a* is simply described as

When creating a node (see 2.2.1), you may want to give some attributes of the node to create. You can give a list of attributes: *type node_name attributes_list* Like the name of a node, most of these attributes are string expressions. The table 2.1 shows which attributes can be used with a type of node.

	place	queue	transition	immediate	type
name	*	*	*	*	string
х	*	*	*	*	arithmetic
у	*	*	*	*	arithmetic
r	*	*	*	*	arithmetic
t	*	*	*	*	arithmetic
domain	*	*			string
marking	*	*			string
guard			*	*	string
priority			*	*	string
delay			*		string
action			*		string
weight				*	string

Table 2.1: Attributes

Therefore, to describe a place named *place* with a marking equals to 1, it's done this way:

```
place "place" marking "1"
or
place "place" (marking "1")
```

If there are more attributes to write, the parentheses are mandatory and attributes are separated by a comma:

```
place "place" (marking "1", domain "red" )
```

As PetriScript supposes that each node has a unique name, it is sufficient to designate it with its name when modifying it, but you can type the whole description if you want to. If one attribute is wrong even if the name is good, then PetriScript will complain about not finding the node.

The attribute **name** is only useful when modifying a node (see 2.2.4). Note that writing **place** "place" (**name** "another name", **marking** "2") is strictly equivalent to **place** "another name"**marking** "2".

The Net node

The Net node is a particular node which has two attributes: **authors** and **declaration**. Only one operation can be made on this special node: the modification. For more details, see subsection 2.2.4, page 16.

Graphical positioning

The x and y attributes are the cartesian coordinates of a node in the Macao window, while r and t are its polar coordinates. When using these latter attributes, PetriScript uses x and y as the center of the orthonormal coordinate system in which polar coordinates are used, as shown in figure 2.1.

The upper left corner of the Macao window is used as the center of the nodes coordinates, the x axis is oriented from the left to the right and the y axis from the top to the bottom.

2.1.2 Lists of nodes

Lists are useful for fusions of sets of nodes. Before using one, it is necessary to declare it as described in 1.6, p. 4.

A such list is a sequence of nodes of the same type separated with a comma and included within '{' and '}'. The type is specified before the list:

Adding nodes

Here are the four ways to add nodes to an existing list:

• variable := type node_list
The content of the variable is lost (but the nodes it formerly contained are not deleted)
and replaced with the given list. If the two given lists are of different types, variable
takes the type of node_list.

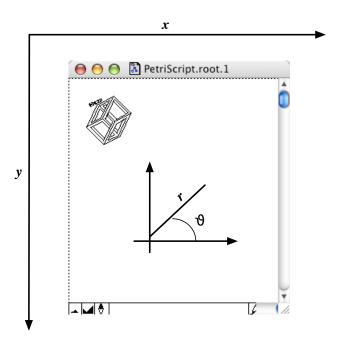


Figure 2.1: Cartesians and polar coordinates in PetriScript

```
$1 := place {"p1", "p2"};
$m := transition {"t1","t2","t3"};
```

• *append* node *to* variable

It simply adds the node *node* at the end of variable *variable*. If *node* and *variable* are of different types, an error is produced.

```
append place "p" to $1;
append transition "t" to $m;
```

• *append* type node_list **to** variable

If *node_list* and *variable* are of different types, an error is produced.

```
append place {"p1","p2"} to $1;
append transition {"t1","t2","t3"};
```

• append variable to variable

If variable and variable are of different types, an error is produced.

```
append $n to $1;
```

Note **1** You can't add a node that doesn't exist to a list. Otherwise, it will produce an error.

Note **2** If you add the same node two or more times to a list, only its first occurrence will be in the list.

Accessing node in list

Accessing to an element in a list is done by giving its position between brackets: *variable*[*position*], where *position* is an arithmetic expression. It is so possible to designate a node via the list containing it. So it is possible to write things like this:

```
append $1[2] to $m;
```

Trying to access an element at a position which is beyond the list's size does nothing and doesn't produce an error. The first element is at position 0 and the last element is at position 1 - 1 for a list 1. A list in an arithmetic expression returns its size.

As an example, the following script creates some places, then adds them to a list and finally prints this list element by element on the historic window.

List inter-dependencies

To prevent the presence of a non-existing node in a list, if a node is contained in more than one list, deleting or modifying it will also delete or modify it in every other list wich has a reference on it.

Regular Expressions

An important feature of PetriScript is that you can add nodes with regular expressions in lists. In fact, anywhere a list is expected, you can use these expressions. These expressions are surrounded by '%' and are the same as in Perl or Python. They return all nodes that match the given expression. For example, if they are three places a, b, c in a Petri net, the list 1 will contain place a and b:

```
list $1;
$1 := place {%[a-b]%};
```

You can use these regular expressions for every attribute: if now places b and c have the same domain d, you can select them with the following script:

```
list $1;
$1 := place {\%.\% domain \%d\%};
```

You can use these every time you need to construct a list, as shown in the following example which adds all places that contain "places" in their names to the list \$1:

Note The regular expression like *ab* will match not only the text *ab*, but also the text *abba*. So if you want to match an exact word, you can use *ab* where 'ab where 'ab

2.1.3 Arcs

They are simply designated this way: (node1 , node2), which is the arc oriented from node1 to node2. For example:

```
(place "p", transition "t")
```

If two nodes of the same type are provided, an error is raised.

2.2 Available operations

For each operation, if it is given a node that doesn't exist, the script fails. Creating a node or modifying a node with the characteristics of an existing one also produces an error.

2.2.1 Creation

create node

```
create place "p";
create place "q" ( domain "red", marking "1");
```

The following example shows how to create a set of places (added to a list) and transitions, using string expressions:

2.2.2 Connection

connect valuation node1 to node2
connect node1 to node2
connect inhibitor valuation node1 to node2
connect inhibitor node1 to node2

As Petri nets are oriented graphs, the resultant arc will be oriented from *node1* to *node2*. To make it inhibited, use the keyword **inhibitor**. *Valuation* is a string expression.

```
connect place "p" to transition "t";
connect transition "t" to place "p";
connect "<x>" place "p" to transition "t";
connect inhibitor transition "t" to place "p";
```

2.2.3 Deletion

delete node delete list delete arc

You can delete a sole node or a list of node. Delete an arc has the effect to disconnect the two given nodes, but it does not delete them.

```
delete place "place";
delete $1;
delete( place "p" , transition "t");
delete $1[$i+1];
```

Remember that

```
delete place "place";
```

is strictly equivalent to

```
delete place "place" (domain "red", marking "1");
```

as PetriScript differentiates nodes only with theirs names.

2.2.4 Modification

set node to attributes_list
set arc to string_expression

You can modify every attributes of a node. To rename a node, use the attribute name.

```
set place "p" to name "q"; set transition "t" to ( name "u", guard "3"); set place "q" to ( x\ 100 , y\ 10 ); set $1[3] to domain "blue";
```

Modifying an arc changes its valuation.

```
set ( place "p" , transition "t" ) to "<y>";
```

It's not possible to change a classic arc to an inhibited one, or the opposite. If you want to, delete the old arc and create a new one with the type you want.

The Net node

To modify the Net node, you would write:

```
set net "net" to (authors "a", declaration "d");
```

But as there is only one net node, it's not mandatory to write its name. So you can simply write:

```
set net to (authors "a", declaration "d");
```

2.2.5 Fusion

Four types of fusion are provided: a node with a node, a list with a node, a list with a list and a list into a single node. Trying to merge a node with itself produces an error. In all cases of fusion, the given node or lists are deleted.

Node with node

```
merge node1 and node2 into attributes_list
merge node1 and node2
```

If no *attributes_list* is given, then PetriScript automatically computes the new name and attributes by concatenating node1 and node2 with an underscore. For example merging a place a with a marking m and a place b gives a place with a name a_b and a marking b.

Examples:

Node with List

```
merge node1 and list into attributes_list
merge node1 and list
merge list and node1 into attributes_list
merge list and node1
```

This is the extension of the previous fusion type. *Node1* is duplicated as many times as there are nodes in *list*. Then each copy of *node1* is merged with a node of *list*. This time, *attributes_list* is used as a pattern: each new resultant node will be created with the characteristics contained in *attributes_list* concatenated with a number starting from 1. If no pattern is given, PetriScript computes one itself.

As an example, the following script creates a place "a" and five places prefixed by "b_" in list \$1. Then it merges this list with the place "a", giving a marking equal to 1 to the newly created nodes.

```
int $i;
list $1;
create place "a";
for $i in 1..5 loop
```

```
create place "b_" & '$i';
end loop;
$1 := place {%b_*%};
merge $1 and place "a" into (marking "1");
```

List with list

merge list1 and list2 into attributes_list merge list1 and list2

Again, this is an extension to the previous fusion type. Now each node of *list1* is duplicated as many times as there are nodes in *list2* and reciprocally. Then each copy of *list1* is merged with the corresponding copy of *list2*, as shown in figure 2.4.

The following script creates two types of places: the ones which are prefixed by "a_" and the ones which are prefixed by "b_". It adds them into lists \$1 and \$m, then merges these lists.

Single list

merge list into attributes_list merge list

Now, all nodes of *list* are merged into a single one, with the attributes of *attributes_list* if presents, or computed automatically if not.

Examples:

```
merge $1 into name "another_place";
merge place {"a1","a2","a3"};
```

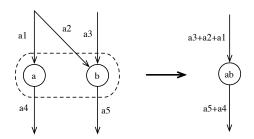


Figure 2.2: Nodes fusion

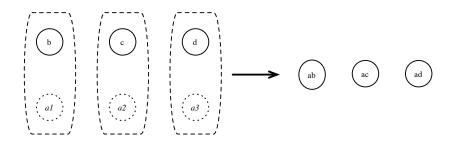


Figure 2.3: Fusion of a node a with a list $\$E = \{b, c, d\}$

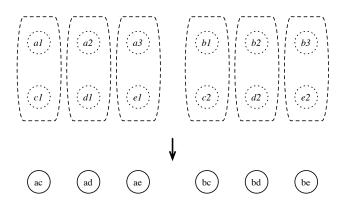


Figure 2.4: Fusion of a list $\$E_1=\{a,b\}$ with a list $\$E_2=\{c,d,e\}$



Figure 2.5: Fusion of a single list $\$E1 = \{a, b, c\}$ into a node c

Chapter 3

Examples

Here are some more realistic examples:

3.1 FIFO

The following script creates a fifo (shown in figure 3.1) using macros and looping instructions.

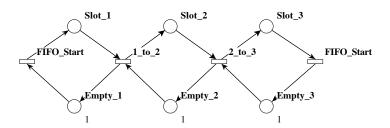


Figure 3.1: FIFO

```
define(FIFO_SIZE,3)
define (FIFO_BASE_X,100)
define (FIFO_BASE_Y,100)
define (FIFO_STEP, 120)
int $wave := 0;
for $wave in 1..FIFO_SIZE loop
       create place "Slot_" & '$wave' (x FIFO_BASE_X + FIFO_STEP * $wave,
       end loop;
for $wave in 1..FIFO_SIZE+1 loop
       create transition "t" & '$wave -1' & "_to_" & '$wave' (x FIFO_BASE_X + FIFO_STEP * $wave - FIFO_STEP / 2,
               y FIFO_BASE_Y + 50);
       if $wave < FIFO_SIZE+1 then
               connect "1" transition "t" &'$wave -1' & "_to_" & '$wave' to place "
                   Slot_" & '$wave';
               connect "1" place "Empty_" & '$wave' to transition "t" &'$wave -1' & "
_to_" & '$wave';
       end if;
       if wave > 1 then
               connect "1" transition "t" &'$wave -1' & "_to_" & '$wave' to place "
                   Empty_" & '$wave - 1';
```

3.2 Philosophers' dinner

This is a coloured version of the Philosophers' dinner, shown in figure 3.2.

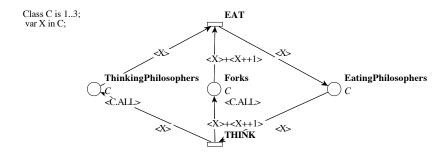


Figure 3.2: Philosophers' dinner

```
define(THK, ThinkingPhilosophers)
define(PHILOS,3)
define(EATING, EatingPhilosophers)
define(XPOS,100)
define(XPOS,100)
define(XINCR,150)
define(XINCR,150)

— declare our variables in the net
set net to declaration "Class C is 1..PHILOS; / var X in C;";

create place "THK" ( domain "C" , marking "<C.ALL>", x XPOS , y YPOS);
create transition "EAT" ( x XPOS + XINCR , y YPOS - YINCR);
create place "EATING" ( domain "C" , x XPOS + 2 × XINCR , y YPOS);
create transition "THINK" ( x XPOS + XINCR , y YPOS + YINCR);
create place "Forks" ( domain "C" , marking "<C.ALL>" , x XPOS + XINCR , y YPOS);

connect "<X>" place "THK" to transition "EAT";
connect "<X>" transition "EAT" to place "EATING";
connect "<X>+<X++1>" place "Forks" to transition "EAT";
connect "<X> place "EATING" to transition "THINK";
connect "<X>+<X++1>" transition "THINK" to place "Forks";
connect "<X>" transition "THINK" to place "Forks";
connect "<X>" transition "THINK" to place "THK";
```

3.3 Another Philosophers' dinner

Much more complex to write, here is the non-coloured version of the philosophers' dinner problem. In fact, the script complexity is due to the graphical positioning of nodes (see figure 3.3).

```
define (X,500)
define (Y,500)
define (RADIUS,80)
```

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```
define (STEP, 40)
define (PHILOS, 3)
define (ANGLE, 360 / PHILOS)
define(ADJUST,70)
int $i := 0;
for $i in 1..PHILOS loop
                              create place "THINK_" & '$i' (marking "1" , x X , y Y, r RADIUS + 6*STEP, t ANGLE
                                              *$i);
                               create transition "TAKE_LEFT_1_FORK_" & '$i' (x X , y Y, r RADIUS + 5*STEP , t
                                             ANGLE*$i-ADJUST/5);
                              create transition "TAKE_RIGHT_1_FORK_" & '$i' (x X , y Y , r RADIUS + 5*STEP , t
                                            ANGLE * $i+ADJUST/5);
                              create place "WAIT_RIGHT_FORK_" & '$i' (x X , y Y, r RADIUS + 4*STEP, t ANGLE*$i-
                                             ADJUST/4);
                              create place "WAIT_LEFT_FORK_" & '$i' (x X , y Y, r RADIUS + 4*STEP, t ANGLE*$i+
                                              ADJUST/4);
                              create transition "TAKE_LEFT_2 FORK_" & '$i' (x X, y Y, r RADIUS + 3*STEP , t
                                          ANGLE* $i + ADJUST / 3);
                              create transition "TAKE_RIGHT_2_FORK_" & '$i' (x X, y Y, r RADIUS + 3*STEP, t
                              ANGLE*$i-ADJUST/3);

create place "EAT_" & '$i' ( x X, y Y, r RADIUS + 2*STEP, t ANGLE*$i);

create transition "RELEASE_FORK_" & '$i' (x X, y Y, r RADIUS + STEP, t ANGLE*$i
                             connect place "THINK_" & '$i' to transition "TAKE_LEFT_1_FORK_" & '$i'; connect place "THINK_" & '$i' to transition "TAKE_RIGHT_1_FORK_" & '$i'; connect transition "TAKE_LEFT_1_FORK_" & '$i' to place "WAIT_RIGHT_FORK_" & '$i'; connect transition "TAKE_RIGHT_1_FORK_" & '$i' to place "WAIT_LEFT_FORK_" & '$i'; connect place "WAIT_RIGHT_FORK_" & '$i' to transition "TAKE_RIGHT_2_FORK_" & '$i'; connect place "WAIT_RIGHT_FORK_" & '$i' to transition "TAKE_RIGHT_2_FORK_" & '$i'; connect place "WAIT_RIGHT_FORK_" & '$i' to transition "TAKE_RIGHT_2_FORK_" & '$i'; connect place "WAIT_RIGHT_FORK_" & '$i' to transition "TAKE_RIGHT_2_FORK_" & '$i'; connect place "WAIT_RIGHT_FORK_" & '$i' to transition "TAKE_RIGHT_2_FORK_" & '$i'; connect place "WAIT_RIGHT_FORK_" & '$i'; connect place "WAIT_RIGHT_5ORK_" & '$i'; connect place "WAIT_FIGHT_5ORK_" & '$i'; connect place "WAIT_FIGHT_5ORK_" & '$i'; connect place "WAIT_FIGHT_5ORK_" & '$i'; connect pl
                             connect place "WAIT_LEFT_FORK_" & '$i' to transition "TAKE_LEFT_2_FORK_" & '$i'; connect transition "TAKE_RIGHT_2_FORK_" & '$i' to place "EAT_" & '$i'; connect transition "TAKE_LEFT_2_FORK_" & '$i' to place "EAT_" & '$i';
                              connect place "EAT_" & '$i' to transition "RELEASE_FORK_" & '$i';
                              create place "FORK_LEFT_" & '$i'( x X, y Y, r RADIUS, t ANGLE*$i - ADJUST);
create place "FORK_RIGHT_" & '$i' (x X, y Y, r RADIUS, t ANGLE*$i + ADJUST);
                             connect transition "RELEASE_FORK_" & '$i' to place "FORK_LEFT_" & '$i'; connect transition "RELEASE_FORK_" & '$i' to place "FORK_RIGHT_" & '$i'; connect transition "RELEASE_FORK_" & '$i' to place "THINK_" & '$i'; connect place "FORK_LEFT_" & '$i' to transition "TAKE_LEFT_1_FORK_" & '$i'; connect place "FORK_RIGHT_" & '$i' to transition "TAKE_RIGHT_1_FORK_" & '$i'; connect place "FORK_LEFT_" & '$i' to transition "TAKE_LEFT_2_FORK_" & '$i'; connect place "FORK_RIGHT_" & '$i' to transition "TAKE_RIGHT_2_FORK_" & '$i'; connect place "FORK_RIGHT_" & '$i' to transition "TAKE_RIGHT_2_FORK_" & '$i'; connect place "FORK_RIGHT_" & '$i' to transition "TAKE_RIGHT_2_FORK_" & '$i'; connect place "FORK_RIGHT_" & '$i' to transition "TAKE_RIGHT_2_FORK_" & '$i'; connect place "FORK_RIGHT_" & '$i' to transition "TAKE_RIGHT_2_FORK_" & '$i'; connect place "FORK_RIGHT_" & '$i' to transition "TAKE_RIGHT_2_FORK_" & '$i'; connect place "FORK_RIGHT_" & '$i' to transition "TAKE_RIGHT_2_FORK_" & '$i'; connect place "FORK_RIGHT_" & '$i' to transition "TAKE_RIGHT_2_FORK_" & '$i'; connect place "FORK_RIGHT_" & '$i' to transition "TAKE_RIGHT_2_FORK_" & '$i'; connect place "FORK_RIGHT_" & '$i' to transition "TAKE_RIGHT_2_FORK_" & '$i'; connect place "FORK_RIGHT_" & '$i' to transition "TAKE_RIGHT_2_FORK_" & '$i'; connect place "FORK_RIGHT_1 & '$i'; connect place "FORK_RIGHT
end loop;
for $i in 1..PHILOS loop
if $i = PHILOS then
                              merge place "FORK_LEFT_" & 'PHILOS' and place "FORK_RIGHT_1" into (name "FORK_" &
                                                      $i', marking "1");
else
                              merge place "FORK_LEFT_" & '\pmi' and place "FORK_RIGHT_" & '\pmi+1' into (name "
                                             FORK_" & '$i', marking "1");
end if:
end loop;
```

3.4 Trains

Two trains circulate in the same direction on a circular railroad, divided in fifteen sections. The two trains can never, for security reasons, be on two contiguous segments. Traffic lights manage the access to each of these sections. The Petri net shown at figure 3.4, model this problem: sections are represented by places Section_1 to Section_15. The presence of a marking into of this place means that a train is present at this location. Traffic lights are modelled by places F1 to F15. The presence of the marking indicate that the light is green, the entry in the section it is guarding is then possible. The passage from a section x to a section y is done at the activation of the transition x to y.

It is a good example of the automatization made possible by PetriScript: it is sufficient to modify the macro SECTIONS to obtain a bigger net, without having to use the graphical interface.

```
define(X,250)
define(Y,350)
define (radius, 50)
define (R, 150)
define (SECTIONS, 15)
define (INNER_ANGLE,360/SECTIONS)
define(OUTTER_ANGLE,360/(2*SECTIONS))
int $i := 0;
int $j := 0;
for $i in 1.. SECTIONS loop
         create place "F" & '$i' ( x X, y Y, r radius, t $i * INNER_ANGLE);
create place "Section_" & '$i' ( x X, y Y, r R, t $i * INNER_ANGLE);
create transition "t" & '$i' & "_to_" & '$i mod SECTIONS + 1' ( x X, y Y, r R, t
               $i * INNER_ANGLE + OUTTER_ANGLE);
end loop;
for $i in 1.. SECTIONS loop
         connect place "Section_" & '$i' to transition "t"&'$i' & "_to_" & '$i mod
              SECTÎONS + 1';
           \textbf{connect transition} \ \ "t" \ \& \ '\$i \ ' \& \ "\_to\_" \ \& \ '\$i \ \textbf{mod} \ SECTIONS \ + \ 1' \ \textbf{to} \ \ \textbf{place} \ \ "Section\_" 
                & '$i mod SECTIONS + 1';
          if i /= 1 then
                   connect place "F" & '$i' to transition "t" & '$i-1' & "_to_" & '$i';
          else
                    connect place "F1" to transition "t" & 'SECTIONS' & "_to_" & '1';
          end if;
          connect transition "t" &'\$i mod SECTIONS + 1' & "_to_" & '(\$i+1) mod SECTIONS +
              1' to place "F" & '$i';
end loop;
for $i in 1.. SECTIONS loop
          if $i \mod 3 = 0 then
                   set place "Section_" & '$i' to marking "1";
                   set place "F" & '$i' to marking "1";
         end if;
end loop;
```

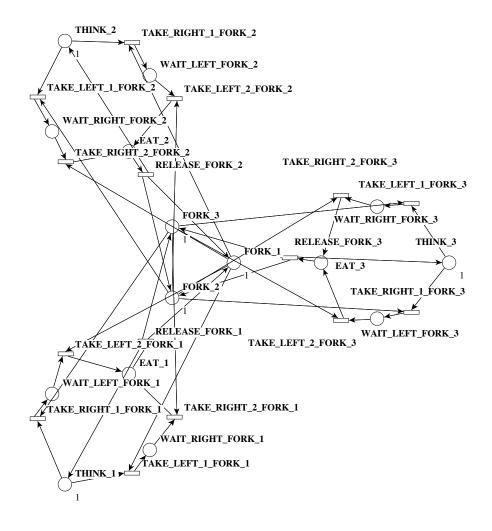


Figure 3.3: Philosopher's dinner

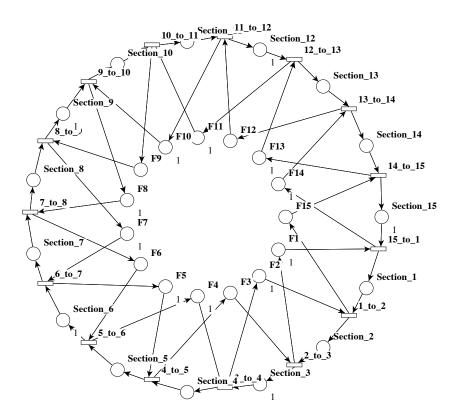


Figure 3.4: Trains

Appendix A

Backus-Naur Form

Keywords are in bold.

A few definitions

```
char ::= ASCII characters, except formatting characters integer ::= [0-9]+ chars ::= sequence of char identifier ::= '$'chars
```

Expressions used by instructions and node descriptions

```
arithmetic\_operator
                                mod
                                '('arithmetic_expression')'
arithmetic_expression
                        1
                                arithmetic\_expression
                                        arithmetic_operator
                                 arithmetic_expression
                                integer
string_expression
                                str_exp_comp
                                str_exp_comp '&' str_exp
                                 "''chars'"'
str_exp_comp
                                 '''identifier'''
                                 '''arithmetic_expression'''
boolean_relator
boolean_operator
                                and
boolean_expression
                                boolean_relation
                                         boolean_operator
```

APPENDIX A. BACKUS-NAUR FORM

```
boolean_relation
                                 not boolean_expression
                                  '(' boolean_expression')'
                                 boolean_relation
boolean\_relation
                                 arithmetic\_expression
                         ::=
                                          boolean_relator
                                 arithmetic\_expression
                                 arithmetic_expression
                                          in
                                 arithmetic_epxression
                                 arithmetic\_expression
                                 '%' chars '%'
regexp
                         ::=
Node description
node_type
                                 place
                                 queue
                                 transition
                                 immediate
                                 net
attribute
                                 name
                                 х
                                 y
                                 domain
                                 marking
                                 guard
                                 priority
delay
                                 action
                                 weight
                                 authors
                                 declaration
attribute\_expression
                                 string_expression
                                 arithmetic_expression
node
                                 node_type node_description
                                 list_access
node_description
                                 string_expression
                                 string_expression attribute attribute_expression
                                 string_expression '('attributes_list')'
attributes\_list
                                 attribute attribute_expression
                         ::=
                                 attribute attribute_expression','attributes_list
List of nodes
                         node_type '{'nodes_list'}'
node_list
                ::=
nodes_list
                         list_component
                         list_component ',' nodes_list
list_component
                         node\_description
                         node_regexp
```

regexp

::=

node_regexp

APPENDIX A. BACKUS-NAUR FORM

```
regexp attribute regexp
                 ı
                         regexp '('list_regexp')'
list_regexp
                 ::=
                         attribute regexp
                         attribute regexp','list_regexp
                1
                         identifier '['arithmetic_expression']'
list_access
                ::=
Arcs
                         '('node, node')'
arc
                 ::=
Main body
body
                         preprocess
                         declarations
                         instructions
                 1
                         instruction ';'
instructions
                ::=
                         instruction ';' instructions
                conditional
instruction
                ::=
                         looping
                 1
                         printing
                 1
                         appending
                         affectation
                         creation
                         connection
                1
                         deletion
                         fusion
                         modification
```

Macros

For more informations on the preprocessor, please take a look at the m4 documentation.

```
preprocess ::= define '('chars','chars')'
```

Declarations

Conditional structures

APPENDIX A. BACKUS-NAUR FORM

Various instructions

Instructions to manipulate nodes

```
creation
                           create node
connection
                           connect inhibitor string_expression node to node
                  ::=
                           connect inhibitor node to node
                           connect string_expression node to node
                  1
                           connect node to node
deletion
                           delete node
                  ::=
                           delete node_list
                           delete identifier
                           delete arc
                           set node to '('attributes_list')'
set net to '('attributes_list')'
set arc to string_expression
modification
fusion
                           merge fusion_comp and fusion_comp into attributes_list
                           merge fusion_comp and fusion_comp
                           merge identifier into attributes_list
                           merge node_list into attributes_list
                           merge identifier
                           merge node_list
                           node
fusion_comp
                           identifier
                           node_list
```

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